Technische Universität München Fakultät für Informatik Lehrstuhl für Effiziente Algorithmen Prof. Dr. Ernst W. Mayr Chris Pinkau

Complexity Theory

Due date: May 25, 2015 before class!

Problem 1 (10 Points)

Prove the existence of a nondeterministic universal Turing machine. That is, prove that there exists a representation scheme of NTMs, and an NTM NU such that for every string α , and input x, $NU(\alpha, x) = M_{\alpha}(x)$:

- 1. Prove that there exists a universal NTM NU such that if M_{α} halts on x within T steps, then NU halts on (α, x) within $cT \log T$ steps, where c is a constant only dependent on α .
- 2. Prove that there is such a universal NTM that runs on these inputs in at most cT steps.

Problem 2 (10 Points)

Show that $\mathbf{SPACE}(n) \neq \mathcal{NP}$. (Note that it is unknown if either class is contained in the other.)

Problem 3 (10 Points)

Define the class $\mathbf{E} = \bigcup_{c} \mathbf{DTIME}(2^{cn})$.

- 1. Is **E** closed under polynomial-time reductions?
- 2. Show that $\mathcal{P}^{\mathbf{E}} = \mathbf{E}\mathbf{X}\mathbf{P}$.

Problem 4 (10 Points)

- 1. Show that \mathcal{NP} and $co\mathcal{NP}$ both are subsets of the set of languages which are polynomial-time Turing reducible to SAT.
- 2. Prove that if \mathcal{NP} was equal to the set of languages which are polynomial-time Turing reducible to SAT, it would follow that $\mathcal{NP} = co\mathcal{NP}$.