6.1 Guessing+Induction

First we need to get rid of the $\mathcal{O}$-notation in our recurrence:

$$T(n) \leq \begin{cases} 2T(\lceil \frac{n}{2} \rceil) + cn & n \geq 2 \\ 0 & \text{otherwise} \end{cases}$$

6.1 Guessing+Induction

First we need to get rid of the $\mathcal{O}$-notation in our recurrence:

$$T(n) \leq \begin{cases} 2T(\lceil \frac{n}{2} \rceil) + cn & n \geq 2 \\ 0 & \text{otherwise} \end{cases}$$

Assume that instead we had

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 2 \\ 0 & \text{otherwise} \end{cases}$$

6.1 Guessing+Induction

First we need to get rid of the $\mathcal{O}$-notation in our recurrence:

$$T(n) \leq \begin{cases} 2T(\lceil \frac{n}{2} \rceil) + cn & n \geq 2 \\ 0 & \text{otherwise} \end{cases}$$

Assume that instead we had

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 2 \\ 0 & \text{otherwise} \end{cases}$$

One way of solving such a recurrence is to guess a solution, and check that it is correct by plugging it in.

6.1 Guessing+Induction

Suppose we guess $T(n) \leq dn \log n$ for a constant d.

6.1 Guessing+Induction

Suppose we guess $T(n) \leq dn \log n$ for a constant d. Then

$$T(n) \leq 2T\left(\frac{n}{2}\right) + cn$$

6.1 Guessing+Induction

Suppose we guess $T(n) \leq dn \log n$ for a constant d. Then

$$T(n) \leq 2T\left(\frac{n}{2}\right) + cn$$
$$\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn$$

6.1 Guessing+Induction

Suppose we guess $T(n) \leq dn \log n$ for a constant d. Then

$$\begin{aligned} T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\ &\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\ &= dn(\log n - 1) + cn \end{aligned}$$

6.1 Guessing+Induction

Suppose we guess $T(n) \leq dn\log n$ for a constant d. Then

$$
\begin{aligned}
T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\
&\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\
&= dn(\log n - 1) + cn \\
&= dn\log n + (c-d)n
\end{aligned}
$$

6.1 Guessing+Induction

Suppose we guess $T(n) \leq dn\log n$ for a constant d. Then

$$\begin{aligned} T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\ &\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\ &= dn(\log n - 1) + cn \\ &= dn\log n + (c-d)n \\ &\leq dn\log n \end{aligned}$$

if we choose $d \geq c$.

6.1 Guessing+Induction

Suppose we guess $T(n) \leq dn \log n$ for a constant d. Then

$$\begin{aligned} T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\ &\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\ &= dn(\log n - 1) + cn \\ &= dn\log n + (c-d)n \\ &\leq dn\log n \end{aligned}$$

if we choose $d \geq c$.

Formally one would make an induction proof, where the above is the induction step. The base case is usually trivial.

6.1 Guessing+Induction

$$T(n) \le \begin{cases} 2T(\frac{n}{2}) + cn & n \ge 16 \\ b & \text{otw.} \end{cases}$$

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$:

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

  Suppose statem. is true for $n' \in \{2, \ldots, n-1\}$, and $n \geq 16$.
  We prove it for n:

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

  Suppose statem. is true for $n' \in \{2, \ldots, n-1\}$, and $n \geq 16$.
  We prove it for n:

$$T(n) \leq 2T\left(\frac{n}{2}\right) + cn$$

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

  Suppose statem. is true for $n' \in \{2, \ldots, n-1\}$, and $n \geq 16$.
  We prove it for n:

$$\begin{aligned} T(n) &\leq 2T\Big(\frac{n}{2}\Big) + cn \\ &\leq 2\Big(d\frac{n}{2}\log\frac{n}{2}\Big) + cn \end{aligned}$$

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

  Suppose statem. is true for $n' \in \{2, \ldots, n-1\}$, and $n \geq 16$. We prove it for n:

$$\begin{aligned} T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\ &\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\ &= dn(\log n - 1) + cn \end{aligned}$$

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

  Suppose statem. is true for $n' \in \{2, \ldots, n-1\}$, and $n \geq 16$.
  We prove it for n:

$$\begin{aligned} T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\ &\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\ &= dn(\log n - 1) + cn \\ &= dn\log n + (c-d)n \end{aligned}$$

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

  Suppose statem. is true for $n' \in \{2, \ldots, n-1\}$, and $n \geq 16$.
  We prove it for n:

$$\begin{aligned} T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\ &\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\ &= dn(\log n - 1) + cn \\ &= dn\log n + (c-d)n \\ &\leq dn\log n \end{aligned}$$

6.1 Guessing+Induction

$$T(n) \leq \begin{cases} 2T(\frac{n}{2}) + cn & n \geq 16 \\ b & \text{otw.} \end{cases}$$

Guess: $T(n) \leq dn \log n$.

Proof. (by induction)

- **base case** $(2 \leq n < 16)$: true if we choose $d \geq b$.
- **induction step** $2 \ldots n-1 \to n$:

Suppose statem. is true for $n' \in \{2, \ldots, n-1\}$, and $n \geq 16$.
We prove it for n:

$$\begin{aligned} T(n) &\leq 2T\left(\frac{n}{2}\right) + cn \\ &\leq 2\left(d\frac{n}{2}\log\frac{n}{2}\right) + cn \\ &= dn(\log n - 1) + cn \\ &= dn\log n + (c-d)n \\ &\leq dn\log n \end{aligned}$$

Hence, statement is true if we choose $d \geq c$.

6.1 Guessing+Induction

Why did we change the recurrence by getting rid of the ceiling?

6.1 Guessing+Induction

Why did we change the recurrence by getting rid of the ceiling?

If we do not do this we instead consider the following recurrence:

$$T(n) \leq \begin{cases} 2T(\lceil \frac{n}{2} \rceil) + cn & n \geq 16 \\ b & \text{otherwise} \end{cases}$$

6.1 Guessing+Induction

Why did we change the recurrence by getting rid of the ceiling?

If we do not do this we instead consider the following recurrence:

$$T(n) \le \begin{cases} 2T(\lceil \frac{n}{2} \rceil) + cn & n \ge 16 \\ b & \text{otherwise} \end{cases}$$

Note that we can do this as for constant-sized inputs the running time is always some constant (b in the above case).

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$T(n)$$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$T(n) \leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn$$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$
\begin{aligned}
T(n) &\leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\leq 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn
\end{aligned}
$$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn\log n$ and get

$$\begin{aligned} T(n) &\leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\ &\leq 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \end{aligned}$$

$$\left\lceil \frac{n}{2} \right\rceil \leq \frac{n}{2} + 1$$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$
\begin{aligned}
T(n) &\leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\leq 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\leq 2(d(n/2+1)\log(n/2+1)) + cn
\end{aligned}
$$

$\left\lceil \frac{n}{2} \right\rceil \leq \frac{n}{2} + 1$

6.1 Guessing+Induction

We also make a guess of $T(n) \le dn\log n$ and get

$$\begin{aligned} T(n) &\le 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\ &\le 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\ &\le 2\big(d(n/2+1)\log(n/2+1)\big) + cn \end{aligned}$$

$\left\lceil \frac{n}{2} \right\rceil \le \frac{n}{2} + 1$

$\frac{n}{2} + 1 \le \frac{9}{16}n$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$
\begin{aligned}
T(n) &\leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\leq 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\
\boxed{\left\lceil \tfrac{n}{2} \right\rceil \leq \tfrac{n}{2} + 1} \quad &\leq 2(d(n/2+1)\log(n/2+1)) + cn \\
\boxed{\tfrac{n}{2} + 1 \leq \tfrac{9}{16}n} \quad &\leq dn \log\left(\frac{9}{16}n\right) + 2d \log n + cn
\end{aligned}
$$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$\begin{aligned} T(n) &\leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\ &\leq 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\ &\leq 2\left(d(n/2+1)\log(n/2+1)\right) + cn \\ &\leq dn \log\left(\frac{9}{16}n\right) + 2d\log n + cn \end{aligned}$$

$\left\lceil \frac{n}{2} \right\rceil \leq \frac{n}{2} + 1$

$\frac{n}{2} + 1 \leq \frac{9}{16}n$

$\log \frac{9}{16}n = \log n + (\log 9 - 4)$

6.1 Guessing+Induction

We also make a guess of $T(n) \le dn \log n$ and get

$$
\begin{aligned}
T(n) &\le 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\le 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\le 2(d(n/2+1)\log(n/2+1)) + cn \\
&\le dn \log\left(\frac{9}{16}n\right) + 2d\log n + cn \\
&= dn \log n + (\log 9 - 4)dn + 2d\log n + cn
\end{aligned}
$$

$\left\lceil \frac{n}{2} \right\rceil \le \frac{n}{2} + 1$

$\frac{n}{2} + 1 \le \frac{9}{16}n$

$\log \frac{9}{16}n = \log n + (\log 9 - 4)$

6.1 Guessing+Induction

We also make a guess of $T(n) \le dn \log n$ and get

$$\begin{aligned}
T(n) &\le 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\le 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\le 2\left(d(n/2+1)\log(n/2+1)\right) + cn \\
&\le dn \log\left(\frac{9}{16}n\right) + 2d\log n + cn \\
&= dn\log n + (\log 9 - 4)dn + 2d\log n + cn
\end{aligned}$$

$\left\lceil \frac{n}{2} \right\rceil \le \frac{n}{2} + 1$

$\frac{n}{2} + 1 \le \frac{9}{16}n$

$\log \frac{9}{16}n = \log n + (\log 9 - 4)$

$\log n \le \frac{n}{4}$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$
\begin{aligned}
T(n) &\leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\leq 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\
\boxed{\left\lceil \tfrac{n}{2} \right\rceil \leq \tfrac{n}{2} + 1} \quad &\leq 2\left(d(n/2 + 1)\log(n/2 + 1)\right) + cn \\
\boxed{\tfrac{n}{2} + 1 \leq \tfrac{9}{16}n} \quad &\leq dn \log\left(\frac{9}{16}n\right) + 2d\log n + cn \\
\boxed{\log \tfrac{9}{16}n = \log n + (\log 9 - 4)} \quad &= dn \log n + (\log 9 - 4)dn + 2d \log n + cn \\
\boxed{\log n \leq \tfrac{n}{4}} \quad &\leq dn \log n + (\log 9 - 3.5)dn + cn
\end{aligned}
$$

6.1 Guessing+Induction

We also make a guess of $T(n) \leq dn \log n$ and get

$$
\begin{aligned}
T(n) &\leq 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\leq 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\
\boxed{\left\lceil \tfrac{n}{2} \right\rceil \leq \tfrac{n}{2} + 1} \quad &\leq 2\big(d(n/2 + 1)\log(n/2 + 1)\big) + cn \\
\boxed{\tfrac{n}{2} + 1 \leq \tfrac{9}{16}n} \quad &\leq dn \log\left(\frac{9}{16}n\right) + 2d \log n + cn \\
\boxed{\log \tfrac{9}{16}n = \log n + (\log 9 - 4)} \quad &= dn \log n + (\log 9 - 4)dn + 2d \log n + cn \\
\boxed{\log n \leq \tfrac{n}{4}} \quad &\leq dn \log n + (\log 9 - 3.5)dn + cn \\
&\leq dn \log n - 0.33dn + cn
\end{aligned}
$$

6.1 Guessing+Induction

We also make a guess of $T(n) \le dn \log n$ and get

$$
\begin{aligned}
T(n) &\le 2T\left(\left\lceil \frac{n}{2} \right\rceil\right) + cn \\
&\le 2\left(d\left\lceil \frac{n}{2} \right\rceil \log \left\lceil \frac{n}{2} \right\rceil\right) + cn \\
\boxed{\left\lceil \tfrac{n}{2} \right\rceil \le \tfrac{n}{2} + 1} \quad &\le 2\left(d(n/2 + 1) \log(n/2 + 1)\right) + cn \\
\boxed{\tfrac{n}{2} + 1 \le \tfrac{9}{16} n} \quad &\le dn \log\left(\frac{9}{16} n\right) + 2d \log n + cn \\
\boxed{\log \tfrac{9}{16} n = \log n + (\log 9 - 4)} \quad &= dn \log n + (\log 9 - 4)dn + 2d \log n + cn \\
\boxed{\log n \le \tfrac{n}{4}} \quad &\le dn \log n + (\log 9 - 3.5)dn + cn \\
&\le dn \log n - 0.33dn + cn \\
&\le dn \log n
\end{aligned}
$$

for a suitable choice of d.