| team | wins | losses | remaining games | | | |
|--------------|-------|----------|-----------------|-----|----|-----|
| i | w_i | ℓ_i | Atl | Phi | NY | Mon |
| Atlanta | 83 | 71 | - | 1 | 6 | 1 |
| Philadelphia | 80 | 79 | 1 | - | 0 | 2 |
| New York | 78 | 78 | 6 | 0 | _ | 0 |
| Montreal | 77 | 82 | 1 | 2 | 0 | - |

Which team can end the season with most wins?

- Montreal is eliminated, since even after winning all remaining games there are only 80 wins.
- But also Philadelphia is eliminated. Why?



Formal definition of the problem:

- Given a set S of teams, and one specific team $z \in S$.
- Team x has already won w_x games.
- Team x still has to play team y, r_{xy} times.
- Does team z still have a chance to finish with the most number of wins.



Flow network for z = 3. *M* is number of wins Team 3 can still obtain.



Idea. Distribute the results of remaining games in such a way that no team gets too many wins.

Certificate of Elimination

Let $T \subseteq S$ be a subset of teams. Define

$$w(T) := \sum_{i \in T} w_i, \qquad r(T) := \sum_{i,j \in T, i < j} r_{ij}$$

wins of
teams in T remaining games
among teams in T

If $\frac{w(T)+r(T)}{|T|} > M$ then one of the teams in T will have more than M wins in the end. A team that can win at most M games is therefore eliminated.



Theorem 1

A team z is eliminated if and only if the flow network for z does not allow a flow of value $\sum_{ij \in S \setminus \{z\}, i < j} r_{ij}$.

Proof (⇐)

- Consider the mincut A in the flow network. Let T be the set of team-nodes in A.
- If for a node x-y not both team-nodes x and y are in T, then x-y ∉ A as otw. the cut would cut an infinite capacity edge.
- We don't find a flow that saturates all source edges:

$$r(S \setminus \{z\}) > \operatorname{cap}(A, V \setminus A)$$

$$\geq \sum_{i < j: \, i \notin T \vee j \notin T} r_{ij} + \sum_{i \in T} (M - w_i)$$

$$\geq r(S \setminus \{z\}) - r(T) + |T|M - w(T)$$

► This gives M < (w(T) + r(T))/|T|, i.e., z is eliminated.

Proof (⇒)

- Suppose we have a flow that saturates all source edges.
- We can assume that this flow is integral.
- For every pairing x-y it defines how many games team x and team y should win.
- The flow leaving the team-node x can be interpreted as the additional number of wins that team x will obtain.
- This is less than $M w_x$ because of capacity constraints.
- Hence, we found a set of results for the remaining games, such that no team obtains more than M wins in total.
- Hence, team *z* is not eliminated.