
Effiziente Algorithmen und Datenstrukturen I

Aufgabe 1 (10 Punkte)

Let $U = \{0, \dots, p-1\}$ for a prime p . For $x \in \mathbb{Z}_p$, define the hash function $h_{a,b}(x)$ as

$$h_{a,b}(x) = (ax + b \pmod p) \pmod n$$

Consider the class of hash functions

$$\mathcal{H} = \{h_{a,b} \mid a, b \in \mathbb{Z}_p\}$$

- Show that \mathcal{H} is not universal.
- Show that \mathcal{H} is (1.1, 2) independent for $p \geq 13$.
- Why would you not choose \mathcal{H} as a class of hash functions?

Aufgabe 2 (10 Punkte)

In double hashing, if we use the hash function $h(k, i) = (h_1(k) + ih_2(k)) \pmod m$, show that when m and $h_2(k)$ have greatest common divisor $d \geq 1$ for some key k , then an unsuccessful search for key k examines $\frac{1}{d}$ th of the hash table before returning to slot $h_1(k)$.

(Note: When $d = 1$, i.e. when m and $h_2(k)$ are relatively prime, the search may examine the entire hash table.)

Aufgabe 3 (10 Punkte)

The mean M of a set of k integers $\{x_1, x_2, \dots, x_k\}$ is defined as

$$M = \frac{1}{k} \sum_{i=1}^k x_i.$$

We want to maintain a data structure D on a set of integers under the normal INIT, INSERT, DELETE, and FIND operations, as well as a MEAN operation, defined as follows:

- INIT(D): Create an empty structure D .
- INSERT(D, x): Insert x in D .
- DELETE(D, x): Delete x from D .
- FIND(D, x): Return pointer to x in D .

5. $\text{MEAN}(D, a, b)$: Return the mean of the set consisting of elements x in D with $a \leq x \leq b$.

Describe how to modify a standard red-black tree in order to implement D , such that INIT is supported in $O(1)$ time and INSERT , DELETE , FIND , and MEAN are supported in $O(\log n)$ time.