

## 6.3 The Characteristic Polynomial

Consider the recurrence relation:

$$c_0T(n) + c_1T(n-1) + c_2T(n-2) + \dots + c_kT(n-k) = f(n)$$

This is the general form of a linear recurrence relation of order  $k$  with constant coefficients ( $c_0, c_k \neq 0$ ).

- The value of  $T(n)$  only depends on the  $k$  preceding values. This means that the order of the recurrence relation is of order  $k$ .
- The recurrence is linear as there are no products of  $T(n)$ 's.
- The coefficients are constant, i.e. they do not change with  $n$ .
- The recurrence relation is homogeneous if  $f(n) = 0$ .
- The recurrence relation is inhomogeneous if  $f(n) \neq 0$ .

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- ▶ The recurrence is linear as there are no products of  $T[n]$ 's.
- ▶ If  $f(n) = 0$  then the recurrence relation becomes a linear, homogenous recurrence relation of order  $k$ .

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### Observations:

- ▶ The solution  $T[0], T[1], T[2], \dots$  is completely determined by a set of boundary conditions that specify values for  $T[0], \dots, T[k-1]$ .
- ▶ In fact, any  $k$  consecutive values completely determine the solution.
- ▶  $k$  non-consecutive values might not be an appropriate set of boundary conditions (depends on the problem).

### Approach:

- ▶ First determine all solutions that satisfy recurrence relation.
- ▶ Then pick the right one by analyzing boundary conditions.
- ▶ First consider the homogenous case.

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# The Homogenous Case

The solution space

$$S = \left\{ T = T[0], T[1], T[2], \dots \mid T \text{ fulfills recurrence relation} \right\}$$

is a **vector space**. This means that if  $T_1, T_2 \in S$ , then also  $\alpha T_1 + \beta T_2 \in S$ , for arbitrary constants  $\alpha, \beta$ .

How do we find a non-trivial solution?

We guess that the solution is of the form  $\lambda^n$ ,  $\lambda \neq 0$ , and see what happens. In order for this guess to fulfill the recurrence we need

$$c_0 \lambda^n + c_1 \lambda^{n-1} + c_2 \cdot \lambda^{n-2} + \dots + c_k \cdot \lambda^{n-k} = 0$$

for all  $n \geq k$ .

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Dividing by  $\lambda^{n-k}$  gives that all these constraints are identical to

$$c_0\lambda^k + c_1\lambda^{k-1} + c_2 \cdot \lambda^{k-2} + \dots + c_k = 0$$

This means that if  $\lambda_i$  is a root (Nullstelle) of  $P[\lambda]$  then  $T[n] = \lambda_i^n$  is a solution to the recurrence relation.

Let  $\lambda_1, \dots, \lambda_k$  be the  $k$  (complex) roots of  $P[\lambda]$ . Then, because of the vector space property

$$\alpha_1\lambda_1^n + \alpha_2\lambda_2^n + \dots + \alpha_k\lambda_k^n$$

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## Lemma 5

Assume that the characteristic polynomial has  $k$  *distinct* roots  $\lambda_1, \dots, \lambda_k$ . Then *all* solutions to the recurrence relation are of the form

$$\alpha_1 \lambda_1^n + \alpha_2 \lambda_2^n + \dots + \alpha_k \lambda_k^n .$$

Proof.

There is one solution for every possible choice of boundary conditions for  $T[1], \dots, T[k]$ .

We show that the above set of solutions contains one solution for every choice of boundary conditions.

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Proof (cont.).

Suppose I am given boundary conditions  $T[i]$  and I want to see whether I can choose the  $\alpha'_i$ 's such that these conditions are met:

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$$\begin{pmatrix} \lambda_1 & \lambda_2 & \cdots & \lambda_k \\ \lambda_1^2 & \lambda_2^2 & \cdots & \lambda_k^2 \\ & & \vdots & \\ \lambda_1^k & \lambda_2^k & \cdots & \lambda_k^k \end{pmatrix} \begin{pmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_k \end{pmatrix} = \begin{pmatrix} T[1] \\ T[2] \\ \vdots \\ T[k] \end{pmatrix}$$

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We show that the column vectors are linearly independent. Then the above equation has a solution.

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- ▶ **base case** ( $k = 1$ ):

A vector  $(\lambda_i)$ ,  $\lambda_i \neq 0$  is linearly independent.



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- ▶ **induction step** ( $k \rightarrow k + 1$ ):  
assume for contradiction that there exist  $\alpha_i$ 's with

$$\alpha_1 \begin{pmatrix} \lambda_1 \\ \vdots \\ \lambda_1^{k-1} \\ \lambda_1^k \end{pmatrix} + \cdots + \alpha_k \begin{pmatrix} \lambda_k \\ \vdots \\ \lambda_k^{k-1} \\ \lambda_k^k \end{pmatrix} = 0$$

and not all  $\alpha_i = 0$ .



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and not all  $\alpha_i = 0$ . **Then all  $\alpha_i \neq 0$ !**



# The Homogeneous Case

$$\alpha_1 \begin{pmatrix} \lambda_1 \\ \lambda_1^2 \\ \vdots \\ \lambda_1^{k-1} \\ \lambda_1^k \end{pmatrix} + \dots + \alpha_k \begin{pmatrix} \lambda_k \\ \lambda_k^2 \\ \vdots \\ \lambda_k^{k-1} \\ \lambda_k^k \end{pmatrix} = 0$$

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$\lambda_1 v_1 =$  (bottom part of the first vector)  
 $\lambda_k v_k =$  (bottom part of the second vector)

# The Homogeneous Case

$$\alpha_1 \begin{pmatrix} \lambda_1 \\ \lambda_1^2 \\ \vdots \\ \lambda_1^{k-1} \\ \lambda_1^k \end{pmatrix} + \dots + \alpha_k \begin{pmatrix} \lambda_k \\ \lambda_k^2 \\ \vdots \\ \lambda_k^{k-1} \\ \lambda_k^k \end{pmatrix} = 0$$

$\lambda_1 v_1 =$  (left vector)       $\lambda_k v_k =$  (right vector)

This means that

$$\sum_{i=1}^k \alpha_i v_i = 0 \quad \text{and} \quad \sum_{i=1}^k \lambda_i \alpha_i v_i = 0$$

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This means that

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Hence,

$$\sum_{i=1}^{k-1} \alpha_i v_i + \alpha_k v_k = 0 \quad \text{and} \quad -\frac{1}{\lambda_k} \sum_{i=1}^{k-1} \lambda_i \alpha_i v_i = \alpha_k v_k$$

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This gives that

$$\sum_{i=1}^{k-1} \left(1 - \frac{\lambda_i}{\lambda_k}\right) \alpha_i \mathbf{v}_i = \mathbf{0} .$$

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$$\sum_{i=1}^{k-1} \left(1 - \frac{\lambda_i}{\lambda_k}\right) \alpha_i v_i = 0 .$$

This is a contradiction as the  $v_i$ 's are linearly independent because of induction hypothesis.

# The Homogeneous Case

## What happens if the roots are not all distinct?

Suppose we have a root  $\lambda_i$  with multiplicity (Vielfachheit) at least 2. Then not only is  $\lambda_i^n$  a solution to the recurrence but also  $n\lambda_i^n$ .

To see this consider the polynomial

$$P(\lambda)\lambda^{n-k} = c_0\lambda^n + c_1\lambda^{n-1} + c_2\lambda^{n-2} + \dots + c_k\lambda^{n-k}$$

Since  $\lambda_i$  is a root we can write this as  $Q(\lambda)(\lambda - \lambda_i)^2$ . Calculating the derivative gives a polynomial that still has root  $\lambda_i$ .

This means

$$c_0n\lambda_i^{n-1} + c_1(n-1)\lambda_i^{n-2} + \dots + c_k(n-k)\lambda_i^{n-k-1} = 0$$

Hence,

$$\underbrace{c_0n\lambda_i^n}_{T[n]} + \underbrace{c_1(n-1)\lambda_i^{n-1}}_{T[n-1]} + \dots + \underbrace{c_k(n-k)\lambda_i^{n-k}}_{T[n-k]} = 0$$

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## What happens if the roots are not all distinct?

Suppose we have a root  $\lambda_i$  with multiplicity (Vielfachheit) at least 2. Then not only is  $\lambda_i^n$  a solution to the recurrence but also  $n\lambda_i^n$ .

To see this consider the polynomial

$$P(\lambda)\lambda^{n-k} = c_0\lambda^n + c_1\lambda^{n-1} + c_2\lambda^{n-2} + \dots + c_k\lambda^{n-k}$$

Since  $\lambda_i$  is a root we can write this as  $Q(\lambda)(\lambda - \lambda_i)^2$ . Calculating the derivative gives a polynomial that still has root  $\lambda_i$ .

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(after taking the derivative; multiplying with  $\lambda$ ; plugging in  $\lambda_i$ )

Doing this again gives

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We can continue  $j-1$  times.

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# The Homogeneous Case

## Lemma 6

Let  $P[\lambda]$  denote the characteristic polynomial to the recurrence

$$c_0T[n] + c_1T[n-1] + \dots + c_kT[n-k] = 0$$

Let  $\lambda_i$ ,  $i = 1, \dots, m$  be the (complex) roots of  $P[\lambda]$  with multiplicities  $\ell_i$ . Then the general solution to the recurrence is given by

$$T[n] = \sum_{i=1}^m \sum_{j=0}^{\ell_i-1} \alpha_{ij} \cdot (n^j \lambda_i^n) .$$

The full proof is omitted. We have only shown that any choice of  $\alpha_{ij}$ 's is a solution to the recurrence.

## Example: Fibonacci Sequence

$$T[0] = 0$$

$$T[1] = 1$$

$$T[n] = T[n - 1] + T[n - 2] \text{ for } n \geq 2$$

The characteristic polynomial is

$$\lambda^2 - \lambda - 1$$

Finding the roots, gives

$$\lambda_{1/2} = \frac{1}{2} \pm \sqrt{\frac{1}{4} + 1} = \frac{1}{2} (1 \pm \sqrt{5})$$

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Hence, the solution is of the form

$$\alpha \left( \frac{1 + \sqrt{5}}{2} \right)^n + \beta \left( \frac{1 - \sqrt{5}}{2} \right)^n$$

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## Example: Fibonacci Sequence

Hence, the solution is

$$\frac{1}{\sqrt{5}} \left[ \left( \frac{1 + \sqrt{5}}{2} \right)^n - \left( \frac{1 - \sqrt{5}}{2} \right)^n \right]$$

# The Inhomogeneous Case

Consider the recurrence relation:

$$c_0T(n) + c_1T(n - 1) + c_2T(n - 2) + \dots + c_kT(n - k) = f(n)$$

with  $f(n) \neq 0$ .

While we have a fairly general technique for solving **homogeneous**, linear recurrence relations the inhomogeneous case is different.

# The Inhomogeneous Case

The general solution of the recurrence relation is

$$T(n) = T_h(n) + T_p(n) ,$$

where  $T_h$  is **any** solution to the homogeneous equation, and  $T_p$  is **one** particular solution to the inhomogeneous equation.

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# The Inhomogeneous Case

Example:

$$T[n] = T[n - 1] + 1 \quad T[0] = 1$$

Then,

$$T[n - 1] = T[n - 2] + 1 \quad (n \geq 2)$$

Subtracting the first from the second equation gives,

$$T[n] - T[n - 1] = T[n - 1] - T[n - 2] \quad (n \geq 2)$$

or

$$T[n] = 2T[n - 1] - T[n - 2] \quad (n \geq 2)$$

I get a completely determined recurrence if I add  $T[0] = 1$  and  $T[1] = 2$ .

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$T[1] = 2$  gives  $1 + \beta = 2 \Rightarrow \beta = 1$ .

# The Inhomogeneous Case

If  $f(n)$  is a polynomial of degree  $r$  this method can be applied  $r + 1$  times to obtain a homogeneous equation:

$$T[n] = T[n - 1] + n^2$$

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Shift:

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Shift:

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