Computer functions

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We want to calculate the SINE:

- 1. Reducing the interval (\mathbb{R} , Float'Range,...) to $\left[-\frac{\pi}{2},\frac{3\pi}{2}\right]$
- 2. Reducing the interval to $\left[-\frac{\pi}{2}, \frac{\pi}{2}\right]$
- 3. Reducing the interval to $\left[-\frac{\pi}{6}, \frac{\pi}{6}\right]$
- 4. Calculating the sine with ...
 - (a) ... a Taylor Polynomial
 - (b) ... a Chebyshev Polynomial

1. Reducing the interval to $\left[-\frac{\pi}{2}, \frac{3\pi}{2}\right]$

We need a function y = f(x) with $0 \le y < 2\pi$ (one period of sine) and

 $\sin x = \sin y.$

So how can we do that? Why is this possible?

The function is called **entier**. It maps every $n \in \mathbb{R}$ to the biggest integer *i* with $i \leq n$. So our formula is

$$\sin y = \sin(x - 2\pi * \operatorname{entier} \frac{x}{2\pi}) = \sin x.$$

As it is much better to have the interval $\left[-\frac{\pi}{2},\frac{3\pi}{2}\right]$ instead of $[0,2\pi]$, we change the above line to

$$\sin y = \sin(x - 2\pi * \operatorname{entier} \frac{x + \frac{\pi}{2}}{2\pi})$$
$$y = (x - 2\pi * \operatorname{entier} \frac{x + \frac{\pi}{2}}{2\pi}).$$

2. Reducing the interval to $\left[-\frac{\pi}{2}, \frac{\pi}{2}\right]$

Now we have the sine on the interval $\left[-\frac{\pi}{2}, \frac{3\pi}{2}\right]$. The sine is symmetric to $y = \frac{\pi}{2}$ so we can reduce the interval to $\left[-\frac{\pi}{2}, \frac{\pi}{2}\right]$ (next slide).



The formula for that is

$$z := \begin{cases} y & \text{if } -\frac{\pi}{2} \le y \le \frac{\pi}{2} \\ \pi - y & \text{if } \frac{\pi}{2} < y \end{cases}$$

To calculate the sine in in the interval we will use a **Taylor Polynominal**. But before, we will reduce the interval some more.

3. Reducing the interval to $\left[-\frac{\pi}{6}, \frac{\pi}{6}\right]$

From the addition theorems we get the following term:

$$\sin(3u) = 3\sin u - 4\sin^3 u.$$

With z = 3u and some small changes we get:

$$\sin z = 3(1 - \frac{4}{3}\sin^2\frac{z}{3})\sin\frac{z}{3}.$$

This can easily be done with the two addition theorems and the Pythagoras

$$sin(a+b) = sin a cos b + sin b cos a$$

$$cos(a+b) = cos a cos b - sin a sin b$$

$$1 = sin^{2}(x) + cos^{2}(x)$$

$$\sin(3u) = \sin(u + 2u)$$

$$= \sin u \cos 2u + \sin 2u \cos u$$

$$= \sin u (\cos(u + u)) + \sin(u + u) \cos u$$

$$= \sin u (\cos^2 u - \sin^2 u) + (2 \sin u \cos u) \cos u$$

$$= \sin u \cos^2 u - \sin^3 u + 2 \sin u \cos^2 u$$

$$= \sin u (3 \cos^2 u - \sin^2 u)$$

$$= \sin u (3(1 - \sin^2 u) - \sin^2 u)$$

$$= \sin u (3 - 3 \sin^2 u - \sin 2u)$$

 $= 3\sin u - 4\sin 3u$

4. (a) Calculating the sine with a Taylor Polynomial

Now we can calculate the sine with a Taylor Polynomial (for example). This can now be done with very little effort since we are close to 0 (we take $x_0 = 0$, then we get an odd function).

A Taylor Polynomial of degree n is defined

$$T_n(x) = \sum_{k=0}^n \frac{1}{k!} * f^{(k)}(x_0) * (x - x_0)^k$$

$$f(x) = T_n(x) + R_n(x,\xi)$$

$$R_n(x,\xi) = \frac{1}{(n+1)!} * f^{(n+1)}(\xi) * (x - x_0)^{n+1}$$

Since we do this for a calculator we want to have 7 digits accuracy. So $R_n(x,\xi)$ has to be $< 0.5 * 10^{-7}$. Our $x_0 = 0$, since it is the centre of our interval, $f(x) = \sin x$, and we try it for n = 8.

$$T_{n}(x) = \sum_{k=0}^{n} \frac{1}{k!} * f^{(k)}(x_{0}) * (x - x_{0})^{k}$$

$$T_{8}(u) = \sum_{k=0}^{8} \frac{1}{k!} * f^{(k)}(0) * u^{k}$$

$$= \frac{1}{1!} * f^{(1)}(0) * u + \frac{1}{2!} * f^{(2)}(0) * u^{2}$$

$$+ \frac{1}{3!} * f^{(3)}(0) * u^{3} + \frac{1}{4!} * f^{(4)}(0) * u^{4} + \dots$$

$$= 1 * 1 * u + \frac{1}{2!} * 0 * u^{2} + \frac{1}{3!} * (-1) * u^{3} + \frac{1}{4!} * 0 * u^{4} + \dots$$

$$= u - \frac{u^{3}}{3!} + \frac{u^{5}}{5!} - \frac{u^{7}}{7!}$$

 $= u - 0.1666666667u^3 + 0.0083333333333u^5 - 0.0001984126984u^7$

$$R_{n}(x,\xi) = \frac{1}{(n+1)!} * f^{(n+1)}(\xi) * (x-x_{0})^{n+1}.$$

$$|R_{8}(u,\xi)| = |\frac{1}{9!} * \sin^{9}(\xi) * (u)^{9}|$$

$$= \frac{|\cos(\xi)| * |(u)^{9}|}{9!} \le \frac{|u|^{9}}{9!}$$
relative error
$$= \frac{|p(u) - \sin(u)|}{|\sin(u)|}$$

$$= \frac{|R_{8}|}{|\sin(u)|} \le \frac{|R_{8}|}{0.95|u|} \le \frac{1.1|u|^{8}}{9!}$$

$$\le \frac{1.1}{9!} (\frac{\pi}{6})^{8} \approx 0.18 * 10^{-7} < 0.5 * 10^{-7}$$

Now we have to write that algorithm optimized for the computer:

$$\begin{split} \tilde{x} &= x * 0.15915494 \\ \tilde{y} &= \tilde{x} - \operatorname{entier}(\tilde{x} + 0.25) \\ \tilde{z} &= \begin{cases} \tilde{y} &, \text{if} - 0.25 \leq \tilde{y} \leq 0.25 \\ 0.5 - \tilde{y} &, \text{if} \ 0.25 < \tilde{y} \end{cases} \\ v &= \tilde{z}\tilde{z} \\ w &= \tilde{z}(3.32464499 + v(-2.43058747 \\ + v(0.53308748 - v * 0.0556757))) \\ \sin x &= w(1.88988158 - ww) \end{split}$$

We have reduced the sine to 8 multiplications and 7 additions.

4. (b) Calculating the sine with a Chebyshev Polynomial

There are other ways then using a Taylor Polynomial. For example can a Chebyshev Polynomial be used, this is better since with the same number of terms it is more accurate. Since this is a much more complicated thing, we have to learn some more basics.

Optimal polynomial approximations

We want to approximate a function F(x) in the closed interval [a,b] by means of a polynomial of degree $\leq n$.

- 1. the polynomial $P_n(x)$ of degree $\leq n$ for which $\max_{\substack{[a,b]}} |P_n(x) - F(x)|$ is as small as possible, if it is absolute error that we are interested in (*minimax-absolute-error*), or
- 2. the polynomial $P_n(x)$ of degree $\leq n$ for which $\max_{[a,b]} |\frac{P_n(x) F(x)}{F(x)}|$ is as small as possible, if it is relative error that we are interested in (*minimax-relative-error*).

Chebyshev's theorem on polynomial approximations

Let u(x) denote a function continuous in a closed, finite interval [a,b], and let v(x) denote a function continuous and nonzero in [a,b]. Let V_n denote the set of polynomials of degree $\leq n$. There exists a unique polynomial $P_n^*(x)$ in V_n such that

$$\max_{[a,b]} |\frac{P_n^*(x)}{v(x)} - u(x)| = \min_{P_n(x) \in V_n} \max_{[a,b]} |\frac{P_n(x)}{v(x)} - u(x)|.$$

. . .

Let $P_n(x)$ denote a polynominal in V_n . Then $P_n(x)$ is $P_n^*(x)$ if and only if there exist $N \ge n+2$ points in [a,b],

$$x_1^* < x_2^* < x_3^* < \dots < x_N^*$$

such that

. . .

$$\frac{P_n(x_k^*)}{v(x_k^*)} - u(x_k^*) = (-1)^k \mu^*$$

k = 1, 2, 3, ..., N,

where

$$|\mu^*| = \max_{[a,b]} |\frac{P_n(x)}{v(x)} - u(x)|.$$

A proof of this theorem is beyond the scope of this lecture. Listeners wishing to study the proof can find one in Achieser $(1956)^*$.

Two ways of construing the foregoing theorem are of interest to us:

*Achieser, N. I. (1956): *Theory of Approximation*. Ungar, New York. english translation by C. J. Hyman.

1. With v(x) = 1 and u(x) = F(x), the function $\frac{P_n(x)}{v(x)} - u(x)$ becomes the **absolute-error function** $P_n(x) - F(x)$. In this case, the theorem asserts that there exists a unique polynimial $P_n^*(x)$ of degree $\leq n$ that approximates F(x) with **minimal absolute error** in [a,b]. The theorem further asserts that $P_n^*(x)$ is uniquely characterized by the fact that the absolute-error function $P_n*(x) - F(x)$ possesses at least n + 2 extreme points in [a,b] at which it is alternately positive and negative and at which the magnitudes of $P_n^*(x) - F(x)$ are equal.

2. With v(x) = F(x) and u(x) = 1, where now it is assumed that $F(x) \neq 0$ in [a, b], the function $\frac{P_n(x)}{v(x)} - u(x)$ becomes the **relative-error function** $\frac{P_n(x)-F(x)}{F(x)}$. In this case the theorem asserts that there exists a unique polynomial $P_n^*(x)$ of degree $\leq n$ that approximates F(x) with **minimax relative error** in *[a, b]*. This $P_n^*(x)$ is uniquely characterized by the fact that the relative-error function $\frac{P_n^*(x)-F(x)}{F(x)}$ posesses at least n+2 extreme points in [a,b] at which it is alternatively positive and negative and at which the magnitude of $\frac{P_n^*(x) - F(x)}{F(x)}$ are equal.

An argument for which the maximum magnitude of the error function is attained is called *critical point*, of the approximation. A minimax polynomial approximation to a function is specifically associated with an integer n and an approximation interval [a,b]. Generally, there is also a difference between the the function with *minimax-absolute-error* and the one with *minimax-relative-error*.

We want a a function of the degree 8 in the interval $\left[-\frac{\pi}{6}, \frac{\pi}{6}\right]$ and we want minimax-absolute-error.

Remez' method for polynomial approximations

This is one of two methods by E. Ya. Remez and is called *Remez' second method*.

We want to approximate F(x) in [a, b] and want to determine the polynomial $P_n^*(x)$ of the degree $\leq n$ that approximates F(x) with minimax absolute error in [a, b]. Let $P_n^*(x) = a_0^* + a_1^*x + \cdots + a_n^*x^n$. For the sake of simplicity, we assume that $P_n^*(x) - F(x)$ is a standard error function. Then $P_n^*(x) - F(x)$ possesses exactly n + 2 critical points in [a, b], including a and b. Let these be denoted by $x_k^*, k = 1, 2, \ldots, n$ and labelled so that

$$a = x_1^* < x_2^* < x_3^* < \dots < x_{n+2}^* = b.$$

Then we know by chebyshev's theorem that

$$a_0^* + a_1^* x_k^* + \dots + a_n^* (x_k^*)^n - F(x)_k^* = (-1)^k \mu^*,$$

 $k = 1, 2, \dots, n+2,$

where

$$|\mu^*| = \max_{[a,b]} |P_n^*(x) - F(x)|.$$

The objective in this method is to compute iteratively the x_s^* 's, μ 's, and the coefficients of $P_n^*(x)$.

1. Initially select n + 2 numbers $x_k, k = 1, 2, ..., n + 2$, such that

$$a = x_1^* < x_2^* < x_3^* < \dots < x_{n+2}^* = b.$$

2. Compute the coefficients of a polynomial $P_n(x) = a_0 + a_1 x + \cdots + a_n x^n$ and the number μ by solving the system of n + 2 linear equations

$$a_0 + a_1 x_k + \dots + a_n (x_k)^n - (-1)^k \mu = F(x_k),$$

 $k = 1, 2, \dots, n+2,$

for n+2 unknowns a_0, a_1, \ldots, a_n , and μ .

3. Locate the extreme points in [a,b] of the absoluteerror function $P_n(x) - F(x)$. For the sake of simplicity, we assume that there are exactly n + 2 extreme points, including a and b. Let these be labelled $y_k, k = 1, 2, ..., n + 2$, where

$$a = y_1 < y_2 < y_3 < \cdots < y_{n+2} = b.$$

4. Replace x_k with y_k for k = 1, 2, ..., n + 2, and repeat the sequence of steps given above beginning with step (2).

 x_k converges to x_k^* , a_k converges to a_k^* , and μ converges to μ^* . The convergence is quadric. A good algorithm to compute the starting values in step 1 is

$$x_k = \frac{1}{2}\cos\frac{(n-k+2)\pi}{n+1} + \frac{1}{2}(b+a),$$

$$k = 1, 2, \dots, n+2.$$

If F(x) is an even or an odd function and the interval is of the form [-a, a], you can use [0, a] instead.

And now we will have fun using Maple.

Enjoy!